## MATH 4290 Homework Assignment 5 Solutions

• If (X, T) is a topological dynamical system and X is a compact metric space with metric d, prove that the quantity  $d_n$  defined by  $d_n(x, y) := \max_{0 \le i < n} d(T^i x, T^i y)$ is a metric for all  $n \in \mathbb{N}$ .

**Solution:** By definition,  $d_n \ge d$ , and so  $d_n(x,y) = 0 \Longrightarrow d(x,y) = 0 \Longrightarrow x = y$  since d is a metric. Similarly,  $d_n(x,x) = \max_{0 \le i < n} d(T^i x, T^i x) = 0$ . Symmetricity is also easy:

$$d_n(x,y) = \max_{0 \le i < n} d(T^i x, T^i y) = \max_{0 \le i < n} d(T^i y, T^i x) = d_n(y, x),$$

where the second inequality uses the fact that d is a metric.

Finally, we must check the triangle inequality:

$$d_n(x,z) = \max_{0 \le i < n} d(T^i x, T^i z) \le \max_{0 \le i < n} (d(T^i x, T^i y) + d(T^i y, T^i z))$$
$$\le \max_{0 \le i < n} d(T^i x, T^i y) + \max_{0 \le i < n} d(T^i y, T^i z) = d_n(x, y) + d_n(y, z).$$

## • Prove the inequalities in Lemma 2.5.1 of the textbook.

**Solution:** The first inequality is proved in the book. The second inequality is  $\operatorname{span}(n, \epsilon) \leq \operatorname{sep}(n, \epsilon)$ . To see this, consider a maximal  $\epsilon$ -separated set S for  $d_n$ . There cannot exist x with  $d_n$ -distance at least  $\epsilon$  from all points in S, as then  $S \cup \{x\}$  would be a larger  $\epsilon$ -separated set for  $d_n$ , a contradiction to maximality of S. Therefore, every point in X is distance less than  $\epsilon$  from some point of S, i.e. S is  $\epsilon$ -spanning for  $d_n$ . The minimal size of such a set is then less than or equal to  $|S| = \operatorname{sep}(n, \epsilon)$ , and so  $\operatorname{span}(n, \epsilon) \leq \operatorname{sep}(n, \epsilon)$ .

Finally, we must prove the third inequality:  $\operatorname{sep}(n, \epsilon) \leq \operatorname{cov}(n, \epsilon)$ . To see this, consider any  $\epsilon$ -separated set S for  $d_n$  and any cover C of X by sets of  $d_n$ -diameter less than  $\epsilon$ . For every  $C \in C$ , C contains at most one point of S; if it contained  $x, y \in S$ , then  $d_n(x, y) \leq \operatorname{diam}(C) < \epsilon$ , contradicting S being  $\epsilon$ -separated. Since every  $C \in C$  contains at most one point of S,  $|C| \geq |S|$ . Since C and S were arbitrary, this shows that the maximum size of such S is less than or equal to the minimum size of such C, i.e.  $\operatorname{sep}(n, \epsilon) \leq \operatorname{cov}(n, \epsilon)$ .

• For the one-sided full shift  $(\{0,1\}^{\mathbb{N}},\sigma)$ , find, with proof,  $\operatorname{sep}(n,2^{-k})$  and  $\operatorname{span}(n,2^{-k})$  for every  $n \in \mathbb{N}$  and  $k \in \mathbb{N} \cup \{0\}$ . (Reminder: the metric here is  $d((x_n),(y_n)) := 2^{-\max\{n \ge 0 : x_i = y_i \ \forall i \le n\}}$ .)

**Solution:** The following fact is immediate from definition of d and will be useful throughout this problem: for any m,  $d((x_n), (y_n)) \leq 2^{-m}$  iff  $x_i = y_i$  for

 $1 \leq i \leq m$ , and  $d((x_n), (y_n)) < 2^{-m}$  iff  $x_i = y_i$  for  $i \leq m+1$  (this is because d takes only values of the form  $2^{-i}$ .)

Now, suppose that S is  $2^{-k}$ -separated for  $d_n$ . Equivalently, for all  $x \neq y \in S$ , there exists  $0 \leq i < n$  so that  $d(\sigma^i x, \sigma^i y) \geq 2^{-k}$ . Equivalently, there exists  $0 \leq i < n$  and  $1 \leq j \leq k+1$  so that  $x_{i+j} \neq y_{i+j}$ . Equivalently, there exists  $1 \leq m \leq n+k$  so that  $x_m \neq y_m$ . Equivalently, the initial words of length n+k in points in S are all distinct. It's clear that the maximal size of such a set is just the number of such words, i.e.  $\operatorname{sep}(n, 2^{-k}) = |A|^{n+k}$ .

Suppose that T is  $2^{-k}$ -spanning for  $d_n$ . Equivalently, for every  $x \in X$ , there exists  $t \in T$  so that  $d(\sigma^i x, \sigma^i t) < 2^{-k}$  for all  $0 \leq i < n$ . Equivalently, for all  $0 \leq i < n$  and all  $1 \leq j \leq k+1$ ,  $x_{i+j} = t_{i+j}$ . Equivalently,  $x_m = t_m$  for all  $1 \leq m \leq n+k$ . Equivalently, for every  $x \in X$ , there exists  $t \in T$  starting with the same n+k-letter word as x. It's clear that the minimal size of such a set is just the number of such words, i.e.  $\operatorname{span}(n, 2^{-k}) = |A|^{n+k}$ .

• If (X,T) factors onto (Y,S), prove that  $h(X,T) \ge h(Y,S)$ .

**Solution:** Choose any  $\epsilon > 0$ , and define  $\phi$  the factor from (X, T) to (Y, S) and  $d_X, d_Y$  the metrics on X, Y respectively. By definition of continuity, there exists  $\delta > 0$  so that  $d_X(x, x') < \delta \Longrightarrow d_Y(\phi x, \phi x') < \epsilon$ . Then,

$$(d_X)_n(x,x') < \delta \Rightarrow \forall i \in [0,n), d_X(T^ix, T^ix') < \delta \Rightarrow \forall i \in [0,n), d_Y(\phi(T^ix), \phi(T^ix')) < \epsilon$$
$$\Rightarrow \forall i \in [0,n), d_Y(S^i\phi x), \phi(S^i\phi x')) < \epsilon \Rightarrow (d_Y)_n(\phi x, \phi x') < \epsilon.$$

Therefore, if  $\mathcal{C}$  is a cover of X by sets with  $(d_X)_n$ -diameter less than  $\delta$ ,  $\phi(\mathcal{C}) := \{\phi(C) : C \in \mathcal{C}\}$  is a cover of Y (by surjectivity) by sets with  $(d_Y)_n$ -diameter less than  $\epsilon$ . By taking  $\mathcal{C}$  to have minimum cardinality, we see that

$$\operatorname{cov}_Y(n,\epsilon) \le \operatorname{cov}_X(n,\delta)$$

Taking logs, dividing by n, and letting  $n \to \infty$  yields  $h_{\epsilon}(Y,S) \leq h_{\delta}(X,T) \leq h(X,T)$ . Letting  $\epsilon \to 0$  yields  $h(Y,S) \leq h(X,T)$ .

• If (X,T) is an isometry (i.e. d(x,y) = d(Tx,Ty) for all  $x, y \in X$ ), prove that h(X,T) = 0.

**Solution:** Since (X,T) is an isometry,  $d(T^ix,T^iy) = d(x,y)$  for all x, y, i, and so  $d_n = d$  for all d. Therefore, for every  $\epsilon > 0$ ,  $\operatorname{cov}(n,\epsilon) = \operatorname{cov}(1,\epsilon)$  for all n, meaning that

$$h_{\epsilon}(X,T) = \lim_{n \to \infty} \frac{\log \operatorname{cov}(n,\epsilon)}{n} = \lim_{n \to \infty} \frac{\log \operatorname{cov}(1,\epsilon)}{n} = 0.$$
So,  $h(X,T) = \lim_{\epsilon \to 0^+} h_{\epsilon}(X,T) = 0.$